

A Ped c on A go n fo Co p ng n e y dzæ on N e of o ee

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Running head: A d c on A o fo y dzæ on

Key words: y dzæ on n, o c o on n fo

1 Abstract

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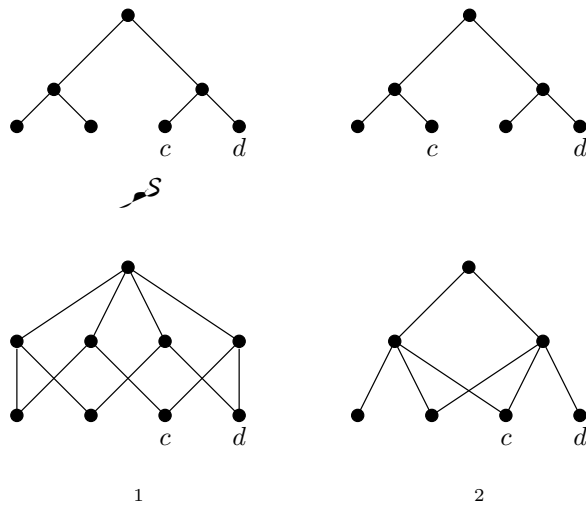


Figure 1. A rooted binary phylogenetic X -tree and a cluster of rooted binary phylogenetic X -trees.

Let X be a finite set of taxa. A rooted binary phylogenetic X -tree is a rooted tree with leaf nodes labeled with elements of X . A cluster of rooted binary phylogenetic X -trees is a set of rooted binary phylogenetic X -trees. A rooted acyclic digraph is a directed graph with no directed cycles. The in-degree of a node is the number of edges entering the node, and the out-degree is the number of edges leaving the node. A hybridization network on X is a rooted acyclic digraph with leaf nodes labeled with elements of X .

3 Reduction Algorithm for Hybridization

Let \mathcal{C} be a cluster of rooted binary phylogenetic X -trees. A rooted binary phylogenetic X -tree is a rooted tree with leaf nodes labeled with elements of X . A cluster of rooted binary phylogenetic X -trees is a set of rooted binary phylogenetic X -trees.

A rooted acyclic digraph is a directed graph with no directed cycles. The in-degree of a node is the number of edges entering the node, and the out-degree is the number of edges leaving the node. A hybridization network on X is a rooted acyclic digraph with leaf nodes labeled with elements of X .

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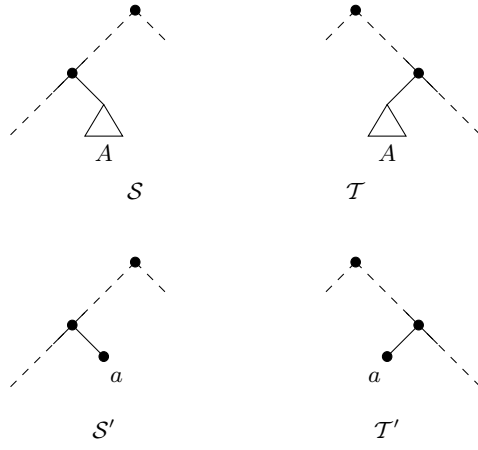


Figure 2. *[The following text is heavily distorted and illegible due to a severe font rendering error.]*

... of Board and ...

Board and ... Hybridization Number ...

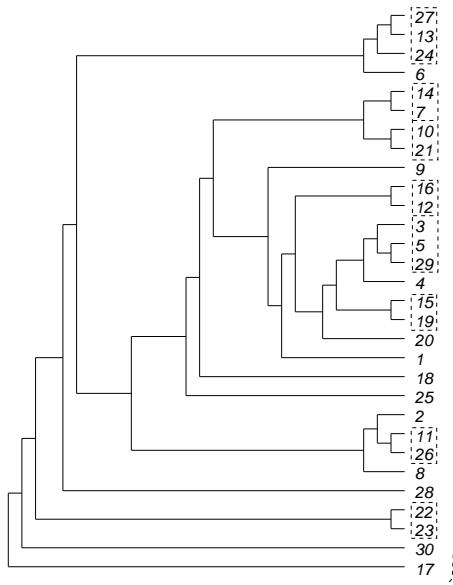
... on nod ... et al. ...

Table 2. Pairwise combinations of Poaceae d

pairwise combination	# taxa	hybridization number	run time ^a
<i>nd</i> - <i>yB</i>	40	14	11 h
<i>nd</i> - <i>cL</i>	36	13	11.8 h
<i>nd</i> - <i>oC</i>	34	12	26.3 h
<i>nd</i> - <i>xy</i>	19	9	320 s
<i>nd</i> - <i>ϕ</i>	46	at least 15	2 d
<i>yB</i> - <i>cL</i>	21	4	1 s
<i>yB</i> - <i>oC</i>	21	7	180 s
<i>yB</i> - <i>xy</i>	14	3	1 s
<i>yB</i> - <i>ϕ</i>	30	8	19 s
<i>cL</i> - <i>oC</i>	26	13	29.5 h
<i>cL</i> - <i>xy</i>	12	7	230 s
<i>cL</i> - <i>ϕ</i>	29	at least 9	2 d
<i>oC</i> - <i>xy</i>	10	1	1 s
<i>oC</i> - <i>ϕ</i>	31	at least 10	2 d
<i>xy</i> - <i>ϕ</i>	15	8	620 s

^arun time on a 2000 MHz CPU, 2 GB RAM machine measured in seconds (s), hours (h), and days (d), respectively

pairwise combinations of *phyB* and *ITS* markers in Poaceae d



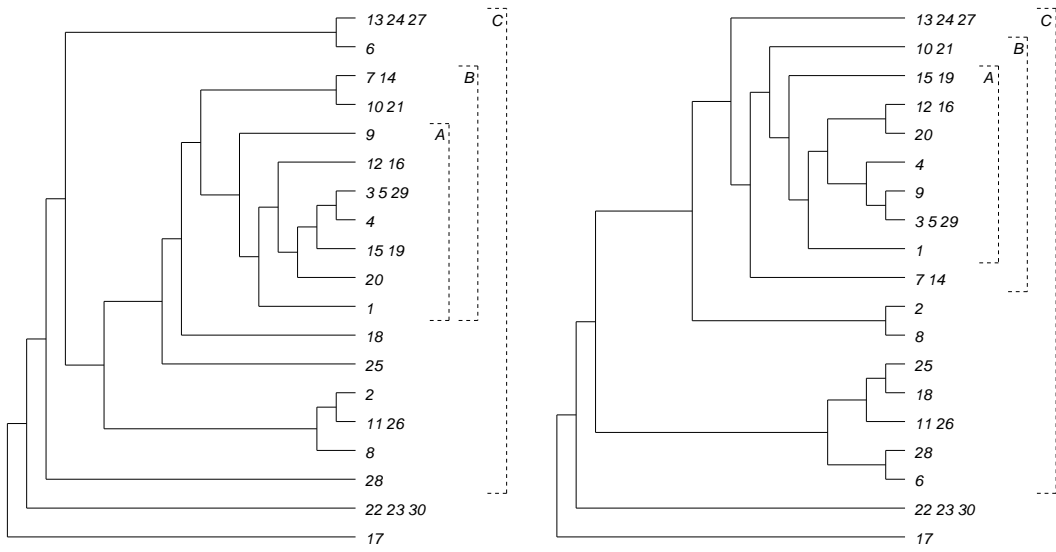
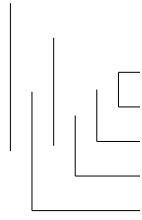


Figure 6.



ny o o c p
nd y d z on n c
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No HybridNumber c c o nd fo n
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o y d z on n n co n z d N n o
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Appendix

A Pseudocode

HybridNumber

```

Algorithm A.1: HybridNumber
  - SubtreeReduction
  - ChainReduction
  if  $n$  consecutive  $C$  of  $S_{nd}$  and
     $\leftarrow C$ 
    do {
      1 - Exhaustive
      2 - ClusterReduction
    }

```

```

Algorithm A.4: ClusterReduction
C ← n
1- C ← c
2- C ← c
1- C ← c
2- C ← c
1- C ← c
2- C ← c
return C

```

```

Algorithm A.5: ExhaustiveSearch
if S return S
repeat
  for each c of S
    do {
      if n cycle n for of S
      do {
        P ← ∑(a,b)∈P
        if P
        do
      }
    }
until
return

```

Remarks

The complexity of the exhaustive search algorithm is $O(2^n)$.
 The complexity of the hybrid number algorithm is $O(n^2)$.
 The complexity of the cluster reduction algorithm is $O(n)$.
 The complexity of the exhaustive search algorithm is $O(2^n)$.
 The complexity of the hybrid number algorithm is $O(n^2)$.
 The complexity of the cluster reduction algorithm is $O(n)$.

